# LINKABILITY IN ITERATED LINE GRAPHS

Thomas  $B\ddot{o}hme^{1)}$ , Martin  $Knor^{2)}$  and  $L'udovít Niepel^{3)}$ 

<sup>1)</sup> Technische Universität Ilmenau, Fak. für Mathematik und Naturwissenschaften, Postfach 100565, 98684 Ilmenau, Deutschland, *E-mail:* tboehme@mathematik.tu-ilmenau.de;

> <sup>2)</sup> Slovak University of Technology, Faculty of Civil Engineering, Department of Mathematics, Radlinského 11, 813 68 Bratislava, Slovakia, *E-mail:* knor@vox.svf.stuba.sk;

 <sup>3)</sup> Kuwait University, Faculty of Science, Department of Mathematics & Computer Science, P.O. box 5969 Safat 13060, Kuwait, *E-mail:* NIEPEL@MATH-1.sci.kuniv.edu.kw.

ABSTRACT. We prove that for every graph H with the minimum degree  $\delta \geq 5$ , the third iterated line graph  $L^3(H)$  of H contains  $K_{\delta \lfloor \sqrt{\delta-1} \rfloor}$  as a minor. Using this fact we prove that if G is a connected graph distinct from a path, then there is a number  $k_G$  such that for every  $i \geq k_G$  the *i*-iterated line graph of G is  $\frac{1}{2}\delta(L^i(G))$ -linked. Since the degree of  $L^i(G)$  is even, the result is best possible.

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### 1. INTRODUCTION AND RESULTS

Let G be a graph. Its line graph L(G) is defined as the graph whose vertices are the edges of G, with two vertices adjacent if and only if the corresponding edges are adjacent in G. Although the line graph operator is one of the most natural ones, only in recent years there is recorded a larger interest in studying iterated line graphs. Iterated line graphs are defined inductively as follows:

$$L^{i}(G) = \begin{cases} G & \text{if } i = 0, \\ L(L^{i-1}(G)) & \text{if } i > 0. \end{cases}$$

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The diameter and radius of iterated line graphs are examined in [10], and [7] is devoted to the centers of these graphs. In [3] and [2], Hartke and Higgins study the growth of the minimum and the maximum degree of iterated line graphs, respectively. The connectivity of iterated line graphs is discussed in [6], and in [13] Xiong and Liu characterize the graphs whose *i*-iterated line graphs are Hamiltonian.

Note that the *i*-iterated line graph of a path on n vertices is a path on n-ivertices for i < n and an empty graph if  $i \geq n$ . The iterated line graph of a cycle is isomorphic to the original cycle, and each iterated line graph of a claw  $K_{1,3}$  is isomorphic to a triangle. Hence, it suffices to study connected graphs distinct from paths, cycles and the claw  $K_{1,3}$ . Such graphs are called *prolific*, since every two members of the sequence  $\{L^i(G)\}_{i=0}^{\infty}$  are non-isomorphic.

Let  $\delta(H)$  denote the minimum degree of H. In [3] we have:

**Theorem A.** Let G be a prolific graph. Then there is  $i_G$  such that for every i,  $i \geq i_G$ , it holds that

$$\delta(L^{i+1}(G)) = 2 \cdot \delta(L^i(G)) - 2.$$

Obviously,  $\delta(L^{i_G}(G)) \geq 3$  in the above theorem. As a consequence, by the results of [6], we obtain:

**Proposition B.** Let G be a prolific graph. Then for every  $i, i \geq i_G + 5$ , the connectivity of  $L^{i}(G)$  equals the minimum degree of  $L^{i}(G)$ .

Here  $i_G$  is the constant appearing in Theorem A.

In this paper we study the linkability of iterated line graphs. A graph with at least 2k vertices is said to be k-linked if for every 2k distinct vertices  $s_1, s_2, \ldots, s_k, t_1$ ,  $t_2, \ldots, t_k$  it contains k vertex-disjoint paths  $P_1, P_2, \ldots, P_k$ , such that  $P_i$  connects  $s_i$  to  $t_i$ ,  $1 \leq i \leq k$ .

Obviously, if a graph is k-linked, then it is k-connected. The converse is far from being true. Jung [4] and, independently, Larman and Mani [8] proved that every 2k-connected graph that contains a subgraph isomorphic to a subdivision of  $K_{3k}$  is k-linked. This together with a result of Mader [9] implies that for every k there is an f(k) such that every f(k)-connected graph is k-linked. Robertson and Seymour [11] extended the result of Jung, Larman and Mani. As a consequence of Theorem (5.4) of [11] we have:

## **Proposition C.** Every 2k-connected graph that has a $K_{3k}$ -minor is k-linked.

In [1] Bollobás and Thomason proved that every 2k-connected graph G with at least 11k|V(G)| edges is k-linked. This implies that every 22k-connected graph is k-linked. Recently, Thomas and Wollan [12] improved the lower bound on the number of edges in the Bollobás and Thomason result to 8k|V(G)|. This was further improved by Kawarabayashi, Kostochka and Yu [5]. They showed that every 2kconnected graph with average degree at least 12k is k-linked. Consequently, every 12k-connected graph is k-linked.

Our main result is the following theorem:

**Theorem 1.** Let G be a prolific graph. Then there is  $k_G$  such that for every  $i \ge k_G$ the graph  $L^i(G)$  is  $\frac{1}{2}\delta(L^i(G))$ -linked.

Observe that a graph with minimum degree  $\delta$  cannot be more than  $\frac{1}{2}\delta$ -linked if  $\delta$  is even. (Consider  $\{s_1, \ldots, s_k, t_1, \ldots, t_k\}$  where  $s_k$  is a vertex of minimum degree 2  $\delta = 2k - 2$ , and  $s_1, \ldots, s_{k-1}, t_1, \ldots, t_{k-1}$  are all of the neighbours of  $s_k$ .) Since the minimum degree of iterated line graph  $L^{i}(G)$  is even if i is "big enough", the result of Theorem 1 is best possible.

We mention that it is an open problem to find to find "good" bounds in terms of G on the numbers  $i_G$  and  $k_G$  in Theorem A and Theorem 1, respectively. However, if the graph G is regular of degree  $\delta$ , then from the proof of Theorem 1 it can be deduced that  $k_G \leq 11$ .

In the proof of Theorem 1, which is trivially true for cycles and the claw  $K_{1,3}$ , we use the following statement:

**Theorem 2.** Let H be a graph with a minimum degree  $\delta \geq 5$ . Then  $L^{3}(H)$  has  $K_t$  as a minor,  $t = \delta \cdot |\sqrt{\delta} - 1|$ .

We remark that the best lower bound for the size of a complete graph in  $L^{3}(H)$ is  $4\delta - 6$ . Theorem 2 shows that there exists a much larger complete graph as a minor.

### 2. Proofs

Let G be a graph and let v be a vertex of  $L^k(G)$ ,  $k \ge 1$ . Then v corresponds to an edge of  $L^{k-1}(G)$ , and this edge will be called 1-history of v. For  $i \geq 2$  we define *i*-histories recursively. The *i*-history of v is a subgraph of  $L^{k-i}(G)$ , edges of which are induced by the vertices of  $L^{k-i+1}(G)$  which are in (i-1)-history of v.

Observe that 1-history is always an edge and 2-history is a path of length 2. The situation is more complicated for *i*-histories when  $i \geq 3$ . The only fact we can say is that *i*-history is a connected graph with at most *i* edges, distinct from any path with less than i edges, see [10]. Therefore we do not visualize the vertices of  $L^{3}(H)$ in H using their 3-histories in the proof of Theorem 2. First we use 2-histories of vertices of  $L^2(H)$  and subsequently 1-histories of vertices of  $L^3(H)$ . In such a way, vertices of  $L^{3}(H)$  correspond to pairs of "adjacent" 2-histories in H.

We prove Theorem 2 in a slightly stronger form. We prove that for an arbitrary vertex v of H there is a subgraph K of  $L^{3}(H)$ , such that  $K_{t}$  is a minor of K and the 3-history of every vertex of K contains v.

## *Proof of Theorem 2.* Denote by $v_1, v_2, \ldots, v_{\delta}, \ldots$ the neighbours of v in H.

Consider 2-histories of the vertices of  $L^2(H)$  in H. Denote by  $c_{i,i'}$  the vertex of  $L^{2}(H)$  with 2-history  $(v_{i}, v, v_{i'})$ , and denote by C the set of these vertices. Then  $|C| \geq {\delta \choose 2}$ . Denote by  $A_i$  those vertices of  $L^2(H)$ , whose 2-history have  $v_i$  as a central vertex and v as an endvertex. Observe that  $|A_i| \geq \delta - 1$ , the vertices of  $A_i$  induce a complete graph in  $L^2(H)$ , and they are adjacent to all  $c_{i,i'}$ ,  $i' \neq i$ . Moreover, the sets  $A_1, A_2, \ldots, A_{\delta}$  are mutually disjoint.

Let  $s = \lfloor \sqrt{\delta - 1} \rfloor$ . Equitably partition every  $A_i$  into s parts  $A_{i,1}, A_{i,2}, \ldots, A_{i,s}$ , so that  $-1 \leq |A_{i,j}| - |A_{i,j'}| \leq 1$  for every  $j \neq j'$ . Then each  $A_{i,j}$  contains at least s vertices, and as  $\delta \geq 5$ , we have  $s \geq 2$ . Denote the vertices of  $A_{i,j}$  by  $a_{i,j,1}, a_{i,j,2}, \ldots, a_{i,j,s}, \ldots$ 

Now denote by  $X_{i,j}$  the set of those vertices of  $L^3(H)$ , whose 1-histories in  $L^2(H)$ contain only the vertices of  $A_{i,j}$ . In the following we show that there are internally vertex-disjoint paths in  $L^{3}(H)$  connecting the sets  $X_{i,j}$ . Let  $X_{i,j}$  and  $X_{i',j'}$  be two such sets,  $(i, j) \neq (i', j')$ . There are two cases to distinguish:

Case 1: i = i'. We join the vertex of  $X_{i,j}$  with 1-history  $(a_{i,j,1}, a_{i,j,2})$  with the vertex of  $X_{i,j'}$  with 1-history  $(a_{i,j',1}, a_{i,j',2})$  by a path of length two. Its interior 3 vertex has 1-history  $(a_{i,j,1}, a_{i,j',1})$ .

Case 2:  $i \neq i'$ . We join the vertex of  $X_{i,j}$  with 1-history  $(a_{i,j,1}, a_{i,j,j'})$  with a vertex of  $X_{i',j'}$  with 1-history  $(a_{i',j',1}, a_{i',j',j})$  by a path of length three. Its interior vertices have 1-histories  $(a_{i,j,j'}, c_{i,i'})$  and  $(c_{i,i'}, a_{i',j',j})$ .

Obviously, the paths just constructed in  $L^3(H)$  are disjoint. If we contract the vertices of  $X_{i,j}$  into a single vertex  $x_{i,j}$ ,  $1 \le i \le \delta$  and  $1 \le j \le s$ , then the vertices  $x_{i,j}$  together with the constructed paths form a subdivision of  $K_{\delta \cdot s}$ . Now the result is a consequence of the fact that all the vertices in  $X_{i,j}$  and in the paths have v in their 3-history.  $\Box$ 

We remark that if  $|A_{i,j}| = s$  in the previous proof, then the paths from  $A_{i,j}$  to  $A_{i',j'}$ , where  $i \neq i'$  and  $j' = 1, 2, \ldots, s$ , exhaust all the vertices with 1-histories  $(a_{i,j,\ldots}, c_{i,i'})$ . This means that the choice  $s = \lfloor \sqrt{\delta-1} \rfloor$  is optimal if we restrict ourselves to the types of paths described in Cases 1 and 2.

Notice that the proof of Theorem 2 implies that, if T is a tree with a central vertex v, such that v and its neighbours have degree  $\delta$  and all the remaining vertices are pendant, then  $L^3(T)$  has  $K_t$  as a minor,  $t = \delta \cdot \lfloor \sqrt{\delta - 1} \rfloor$ .

Proof of Theorem 1. Choose  $k_G$  such that

$$k_G \ge i_G + 5$$
 and  
 $\left\lfloor \sqrt{\delta(L^{k_G - 3}(G)) - 1} \right\rfloor \ge 12,$ 

where  $i_G$  is the constant from Theorem A. Then for every  $i \ge k_G$ , it follows from Proposition B that  $L^i(G)$  is  $\delta(L^i(G))$ -connected. Further, by Theorem A we have  $\delta(L^i(G)) = 8\delta(L^{i-3}(G)) - 14$ . Finally, by Theorem 2  $L^i(G)$  has a  $K_t$ -minor with

$$t = \delta(L^{i-3}(G)) \left\lfloor \sqrt{\delta(L^{i-3}(G)) - 1} \right\rfloor \ge \frac{1}{8} \left( \delta(L^i(G)) + 14 \right) \cdot 12 > \frac{3}{2} \delta(L^i(G)).$$

By Proposition C this implies that  $L^i(G)$  is  $\frac{\delta(L^i(G))}{2}$ -linked.  $\Box$ 

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